## [GTN XLIX:7] MINIMAL CONFIGURATIONS AND INTERLACING

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#### **Abstract**

A graph is singular of nullity  $\eta$  if zero is an eigenvalue of its adjacency matrix with multiplicity  $\eta$ . A subgraph that forces a graph to be singular is called a minimal configuration. We show various properties of minimal configurations.

#### 1. Introduction

Let G = G(V(G), E(G)) be a graph with vertex set  $V(G) = \{v_1, v_2, ..., v_n\}$  and edge set E(G). The adjacency matrix A(G) = A is  $(a_{ij})$ , where  $a_{ij} = 1$  if  $v_i v_j \in E(G)$  and  $a_{ij} = 0$  otherwise. The values of  $\lambda$  for which there exist non-zero vectors  $\mathbf{x}$  such that  $A\mathbf{x} = \lambda \mathbf{x}$  are called eigenvalues of G. The vectors  $\mathbf{x}$  are said to be eigenvectors of G. A graph is singular of nullity G if zero is an eigenvalue of its adjacency matrix G with multiplicity G.

There are subgraphs called *minimal configurations* that are found in singular graphs. We study properties of minimal configurations that determine their structure.

### 2. Singular Graphs

**Definition 2.1:** A kernel eigenvector  $\mathbf{x}_0$  of a singular graph with adjacency matrix  $\mathbf{A}$  is a non-zero vector in the nullspace of  $\mathbf{A}$ .

**Remark 2.2:** A kernel eigenvector  $\mathbf{x}_0 \in \mathbb{R}^n$  of a singular graph G satisfies  $A\mathbf{x}_0 = \mathbf{0}$ .

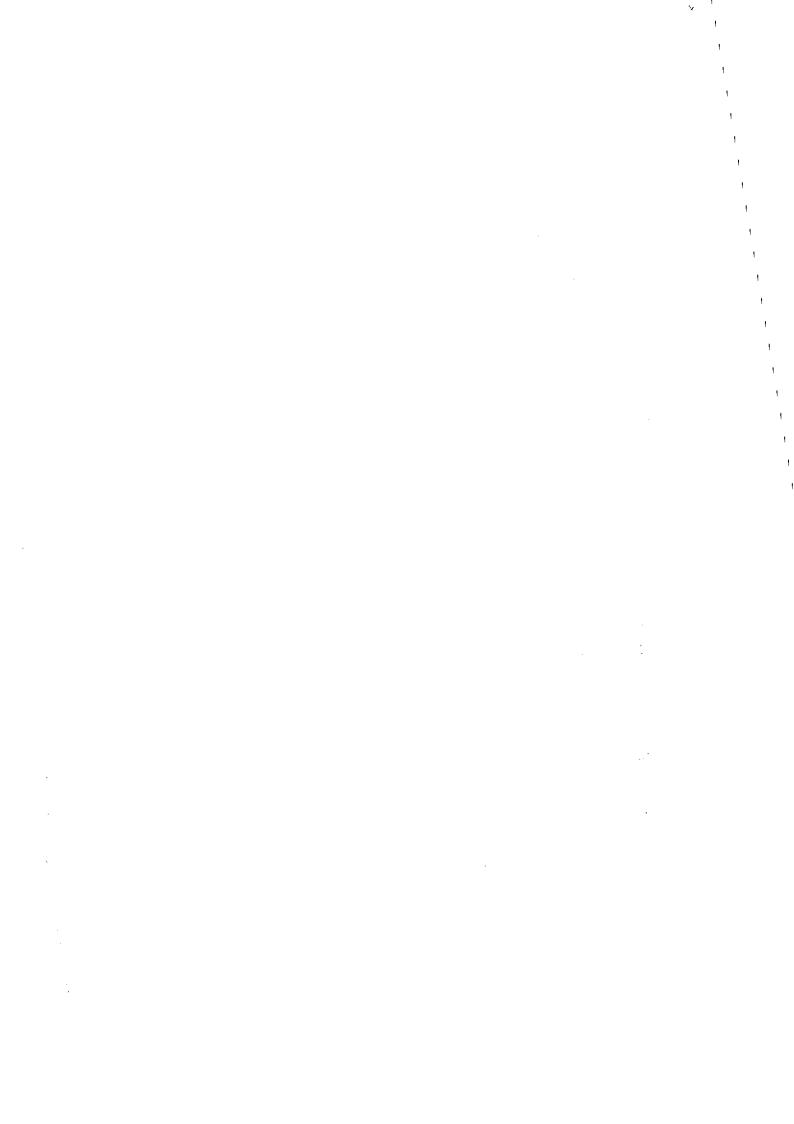
**Definition 2.3:** Let  $x_0$  be a kernel eigenvector of a singular graph G of order  $n \ge 3$ . A subgraph of G induced by the vertices corresponding to the non-zero entries of  $x_0$  is said to be a *core*,  $\chi$ , (with respect to  $x_0$ ). The core is sometimes denoted by  $\chi_p$ , or by  $\chi_{x_0}$ , where p, the number of vertices of the core, is called the *core order*.

The following definition of minimal configurations—the building blocks of singular graphs—is given in [1].

**Definition 2.4:** A singular graph  $\Gamma$  of order  $n \ge 3$ , having a core  $\chi_p$  and *periphery*  $\mathcal{P} := V(\Gamma) - V(\chi_p)$  is said to be a *minimal configuration*, of core order p, if the following conditions are satisfied:

- $(1) \eta(\Gamma) = 1,$
- (2)  $P = \emptyset$  (that is, P is empty, no vertex) or P induces an empty graph (no edge), and
- (3) for  $P \neq \emptyset$ ,  $\eta(\Lambda(\Gamma \setminus v)) > 1$ , for all  $v \in P(\Gamma)$ .

We show that condition (3) in the definition can be replaced by the more powerful condition in the following theorem.



Theorem 2.5: A graph I is a minimal configuration if and only if

- (1)  $\eta(\Gamma) = 1$ ,
- (2)  $P = \emptyset$  or P induces an empty graph, and
- (3)  $\eta(\chi_n) = 1 + |P|$ .

**Proof:** By the *interlacing theorem*, addition to or removal of a vertex, from a graph, changes the nullity by at most one.

Let G be a minimal configuration, then  $\eta(G) = 1$ . By (3) of Definition 2.4, and the interlacing theorem,  $\eta(G \setminus v) = 2$  if  $v \in \mathcal{P}$ .

<u>Case 1</u>:  $P \neq \emptyset$ . In this case there are zero entries in  $x_0$ , the kernel eigenvector of G, which may be written

$$\mathbf{x}_0 = (\alpha_1, \alpha_2, ..., \alpha_{|\gamma|}, 0_1, 0_2, ..., 0_{|\mathcal{P}|}),$$

where  $\alpha_i, \ 1 \le i \le |\chi|$  are the only non-zero entries of  $x_0$ . There is a linear combination, Rel, of the rows or columns of A with coefficients equal to the entries of  $x_0$ . Deletion of a vertex  $v_j \in P$  leaves Rel unaffected since the row or column corresponding to  $v_j$  is not involved. This means that although Rel remains valid, by (3) of Definition 2.4, Rel corresponds to two linearly independent kernel eigenvectors on deletion of any vertex of P. Since no two vertices of P are adjacent,  $v_1, v_2, \ldots, v_{|P|}$  can be successively removed from P0 in any order, increasing the nullity by one with each deletion, until the core P1 is obtained, so that P1 is

<u>Case 2</u>:  $\mathcal{P} = \emptyset$ . In this case  $\Gamma = \chi$  and  $\eta(\chi) = 1$ ;  $\Gamma$  is said to be a *nut graph* [2].

In both cases,  $\eta(\chi) = 1 + |T|$ , as required.

Conversely, let  $\eta(\chi) = 1 + |T|$  and  $\eta(\chi + \nu_1 + \nu_2 + ... + \nu_{|T|}) = \eta(G) = 1$ . Thus,  $\eta(\chi) \ge 1$ . If  $\eta(\chi) = 1$ , then  $|T| = \emptyset$  and  $\chi = G$  is a net graph. If  $\eta(\chi) > 1$ , then  $|T| \ge 1$ . By the Interlacing theorem,

$$|\mathcal{P}| - (j-1) \le \eta(\chi + v_1 + v_2 + \dots + v_j) \le |\mathcal{P}| + j + 1.$$

For j=|T|,  $1 \le \eta(\chi+v_1+v_2+\ldots+v_{|T|}) \le 2|T|+1$ , independently of the sequence in which the vertices of T are added to  $\chi$ . However, the nullity of G is one. This forces the nullity to decrease by unity with each addition of a vertex to  $\chi$ . In particular,  $\eta(G|v_j)=2$  for all  $v_j \in T$  as required.

**Theorem 2.6:** The deletion of a vertex of the core,  $\chi$ , from a minimal configuration  $\Gamma$  produces a non-singular graph.

**Proof:** Let  $\Gamma$  be a minimal configuration with kernel eigenvector  $x_0$  and let  $\nu \in \chi$ , the core of  $\Gamma$ . Suppose  $\Gamma \setminus \nu$  is singular, then  $\Gamma$  has another kernel eigenvector with a zero entry corresponding to  $\nu$ , so that  $\eta(\Gamma) > 1$ , a contradiction.

Remark 2.7: A minimal configuration, with kernel eigenvector  $\mathbf{x}_0$ , corresponding to a core  $\chi_{\mathbf{x}_0}$ , can be considered to be a graph of mullity one with a minimal number of edges and vertices. The labelled vertices in the Figure refer to the vertices of the core  $\overline{K}_4$  in each minimal configuration.

 $\overline{K}_4$  in each minimal configuration  $\Gamma$  is a subgraph of a graph G such that the non-zero entries of the kernel eigenvector are preserved, then G is singular.

For G to be singular with minimal configuration  $\Gamma$ , it is sufficient that the core vertices of  $\Gamma$  have the same degrees as they do in G.

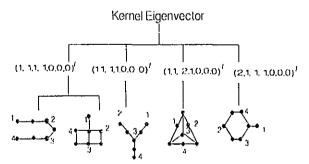


Figure: Minimal configurations with core  $\overline{K}_4$ .

Corollary 2.8: A minimal configuration is connected.

Proof: Suppose that a minimal configuration  $\Gamma$  has at least two components. Since  $\eta(\Gamma)=1$ , only one component, H, is singular and  $V(\chi) \subset V(H)$ . The vertices in the periphery P induce an empty graph  $\Gamma VV(\chi)$ . However,  $\Gamma$  then has isolated vertices and  $\eta(\Gamma)>1$ , a contradiction.

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## References

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